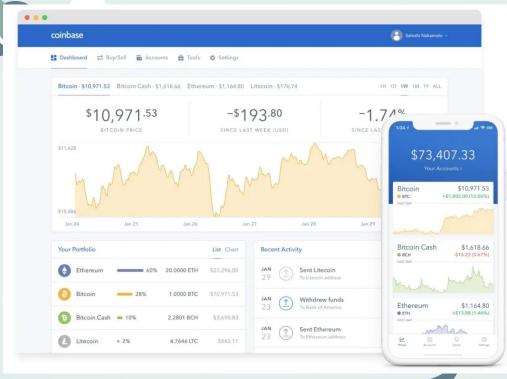
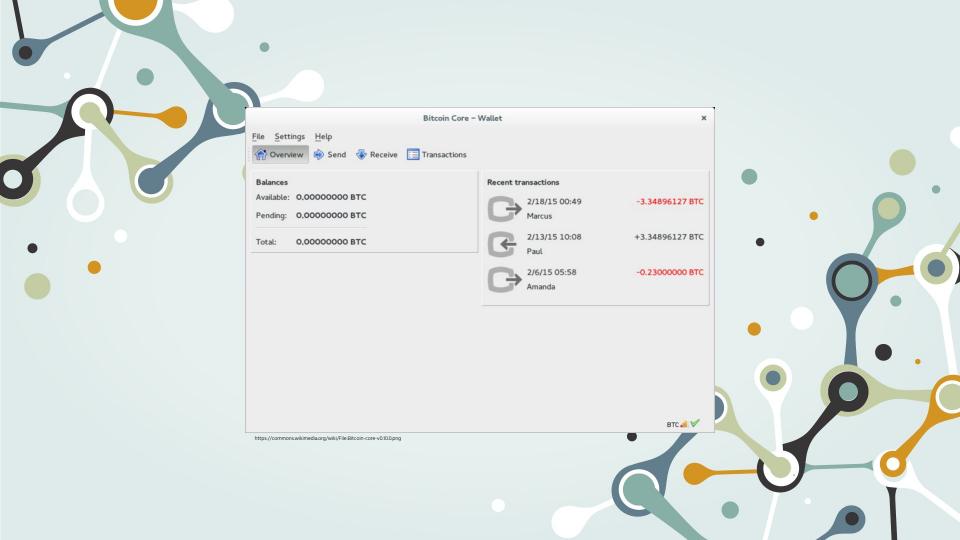
# Elliptic Curve Cryptography and Its Application in Bitcoin Wallet

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https://www.computerworld.com/article/3389678/whats-a-crypto-wallet-and-does-it-manage-digital-currency.html





## What is Bitcoin?

It is a decentralized digital currency that is independent of banks and can be sent from user to user on the peer-to-peer bitcoin blockchain network without the need for intermediaries.



# The Basics

## Private key

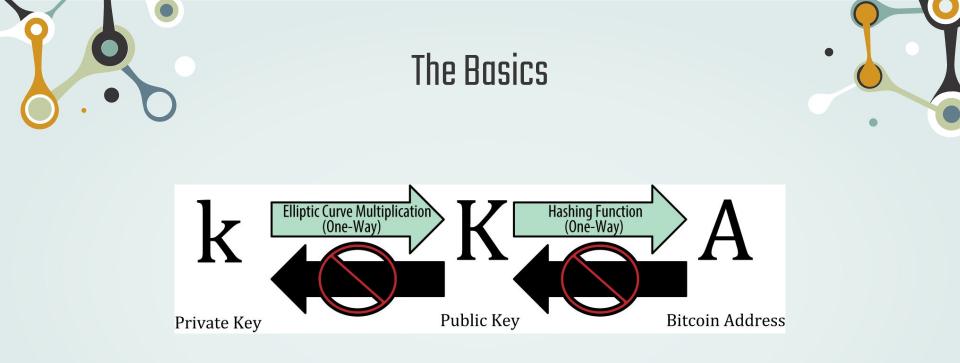
It is the password to your bitcoin account. You do not want to share it with anyone!



### Public key

It is used to prove that the digital signature came from the private key. The signature proves ownership of the private key.

## Address

The bitcoin address is visible to everyone on the blockchain. Just like your bank account number, the sender needs that to transfer money to you. 

Trapdoor functions

# **JUMP IN DEPTH**

How do we obtain those three things?





## How do we create a private key (k)?

Hex: DD5113FEDED638E5500E65779613BDD3BDDBEB8EB5D86CDD3370E629B02E92CD

Base64: 3VET/t7WOOVQDmV3lhO9073b64612GzdM3DmKbAuks0=

WIF: 5KVkpWGfDQGJAUEEDUFbrFxwNPjmXy5kBBmRzzBDf4JkgFXqXTa

#### **Binary:**

https://www.freecodecamp.org/news/how-to-generate-your-very-own-bitcoin-private-key-7ad0f4936e6c/

Bitcoin uses the SHA-256 hash algorithm to generate a secret number k that is 256 bits long.

#### Would there be duplicate private keys?

We don't need to concern about that right now because  $2^{256}$  or  $10^{77}$  is a very big integer.



## How do we create a public key(p)?

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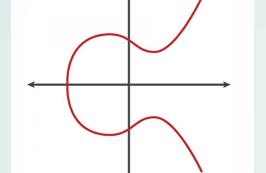


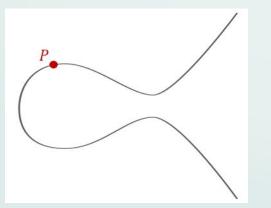
Elliptic Curve Digital Signature Algorithm

https://metamug.com/article/security/sign-verify-digital-signature-ecdsa-java.html

## ECC algorithm is the base of ECDSA

(Elliptic Curve Cryptography)





The function:  $y^2 = x^3 + ax + b$  (a=0,b=7 in bitcoin blockchain)

Let's say:

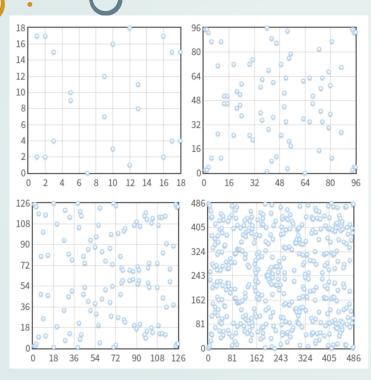
- k is private key
- p is public key
- G(x,y) is a point on the curve, we call it the "generator"
- '\*' is a group operation such that k \* G = (kx, ky)= p (In other words, p is a point)

k \* G = G+G+...+G

Point and point addition: A + B -> C

What if it is A + A? A + A -> use tangent line

## The curve is actually a finite field!



https://medium.com/@\_ericprice/bitcoins-signing-algorithm-elliptic-curves-with-finite-fields-e386f8d0c05

#### How do we pick a G?

- Any point in the block can be used as a generator.
- Everyone uses the same G in bitcoin:  $p=2^{256}-2^{32}-2^9-2^8-2^7-2^6-2^4-2^0$

=0279be667ef9dcbbac55a06295ce870b07029bfcdb 2dce28d959f2815b16f81798

#### How many G are there in the block?

- Group size - 1 since it is a cyclic group.

#### Why using a field?

 To avoid numerical precision issues, we use finite fields because the number of elements is finite and known exactly.



## An example: $y^2 = x^3 + 2x + 3 \mod 97$ With G=(3,6)

0\*G = (infinity) 1\*G = (3,6) 2\*G = (80,10) 3\*G = (80,87) 4\*G = (3,91) 5\*G = (infinity) 6\*G = (3,6) 7\*G = ... The order of this cyclic group is 5, much less than 97.

The points have been partitioned into separate cyclic groups, all of the same size.

# of such groups is call
"cofactor", h=n/r





## How fast can we compute the public key p

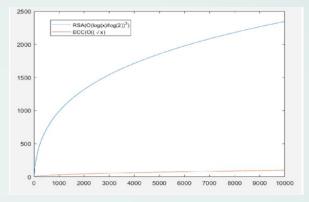
Do we have to do calculations k times?

 No! The time complexity of ECC algo is O(log<sub>2</sub>n)

```
why?
G+G=2G, 2G+2G=4G, 4G+4G=8G
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•••
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If n = 16, we need to do sqrt(16)=4 calculations.

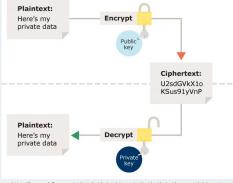


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https://www.researchgate.net/figure/The-time-complexity-of-RSA-and-ECC_fig4_33083214
```

We love ECC because it is much faster than RSA with the same key size.

(1) Data Encryption: If we know p, we can hardly compute k.

Alice: 
$$k_{a}, p_{a} \rightarrow k_{a} p_{b} = G k_{a} k_{b}$$
  
Bob:  $k_{b}, p_{b} \rightarrow k_{b} p_{a} = G k_{a} k_{b}$ 

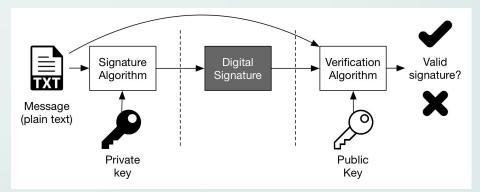


https://ico.org.uk/for-organisations/guide-to-data-protection/guide-to-the-general-data-prote ction-regulation-gdpr/encryption/what-types-of-encryption-are-there/

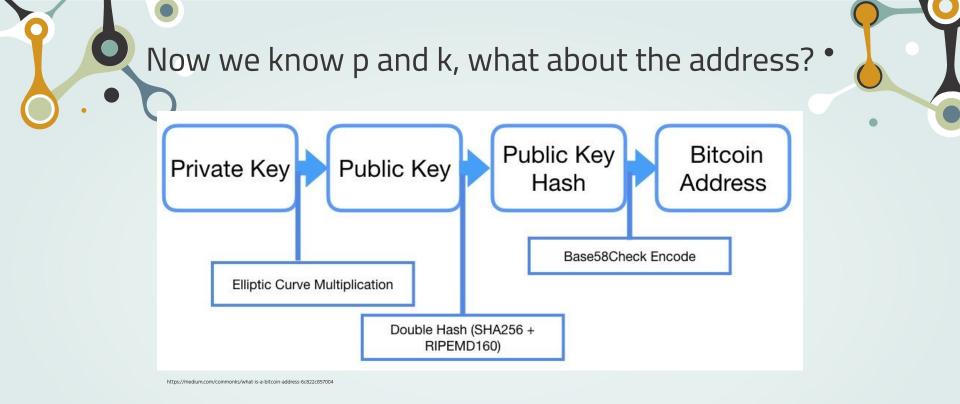
#### (2) Digital Signature: If we are given k, we can compute whether k\*G matches p.

Suppose m is message m\*k = N(digital signature) m\*p = m\*G\*k = (m\*k)\*G=N\*G

We can check whether signature matches to verify the integrity of the message m.



https://stakey.club/en/verifying-digital-signatures/



## We can compute the Bitcoin address from Double Hash functions(SHA256 + RIPEMD160)



## **Issues With ECC Implementation**



#### What if k is a static value? Disaster!

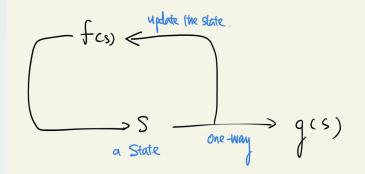
A secure implementation of the ECC curve is theoretically possible, it is not easy to achieve. There are numerous examples of how failed implementation of ECC algorithms resulted in significant vulnerabilities in the cryptographic software.

A great example is that of the Sony ECDSA security disaster. Although Sony used ECDSA to sign software for their PlayStation game console, they did not properly implement the algorithm. Using static parameters instead of random ones made Sony's implementation of the algorithm solvable and subsequently useless.



## Can we obtain truly randomness?

AKA can we build a random number generator that guarantees to create a random number?



Typical workflow of a random number generator







# Thank you!